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Privacy-friendly Platform for Healthcare Dat. ir Cloud Based on Blockchain Environment

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Abstract

Data in cloud has always been a point of attraction for the cyber attackers. Nowadays healthcare data in cloud has become their new interest. Attacks on these healthcare data can result in annihilating consequences for the healthcare organizations. Decentralization of fiese cloud data can minimize the effect of attacks. Storing and running computation on sensitive private healthcare data in cloud are possible by decentralization which is enabled by peer to peer (P2P) network. By leveraging the decentralized or distributed property, blockchain technology ensures the accountability and integrity. Different solutions have been proposed to control the effect of attacks using decentralized approach but these solutions

 $^{^{&}lt;table-row>}$ A preliminary version of this paper appears in The 10th International Conference on Security, Privacy and Anon mity $^{}$ m Computation, Communication and Storage, SpaCCS Workshops 2017. This is the full version.

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somehow failed to ensure overall privacy of patient centric system. Ir this paper, we present a patient centric healthcare data management system wing which helps to attain privacy. Cryptographic tunctions are used to encrypt patients data and to ensure pseudonymity. We manyze the data processing procedures and also the cost effectiveness of the sname contracts used in our system.

Keywords: Blockchain, Decentralization, Healthcare data in cloud, Pseudonymity, Privacy, Security, Smart contract

1. Introduction

A lot of work is going on healthcare and incompation technology in an amalgamated manner and these works are bring a lot of changes in healthcare discipline. These changes are affecting patients attement process hence requiring careful data processing. For treatment, healthcare is completely dependent on data which arises some concerns over data courity and privacy. Authorization or private access to the personal data of introdual patient refers to the term Privacy, which means only authenticated parties will be able to access the private data. Keeping these personal data safe from the eavesdroppers or intruders refers to the term Security, which means system will be able to protect users' private data from outsiders. Authenticated parties of realthcare data preservation process will get

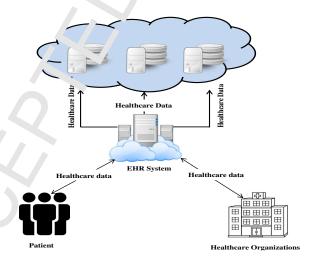


Figure 1: Entities of EHR system and it's Data flow

the access to store data into cloud and retrieve from it. Interaction between the system and the patient requires a secured channel. Different authorication protocol [20, 19, 24] have been proposed to preserve the privacy ar a security. Lack of security may result in devastating consequences like data los and data theft. A lot of intruders are searching for an insecured channel and trying access valuable healthcare data in the cloud network. In most of the case, data 'oss in healthcare causes detrimental consequences to the patients and health care organizations. Due to recent attacks on healthcare data in cloud systems, anferent countries like USA [8] and UK [12] have experienced critical data los. Fersonal data of patients' were kept without encryption in the cloud which allowed "Le attackers to steal the sensitive private data. Let's assume a scenario where patients keep their data in any Electronic Health Record (EHR) system [35, 14, 13, 38, 5] for preservation and also for further access. Figure 1 depicts a generalized formation of EHR systems. In the figure patients and healthcare of mulzations take part in the process as both data sender and data receiver. EHR system is the manager of the whole process that maintains the data flow of the sy, 'er . 10p most entity is the cloud where data is kept. Patients share their perso all data with the doctors and healthcare organizations with the help of these EHR systems. Suppose, a patient keeps her data in the cloud system [7] which uses works ain as a data storage platform. System will store the data on blockchain when in a patient shares her data with the system. Accountability of data is system centric in case of the instance [7], whereby the system will provide data stor ge serv ce even when data is shared with the doctors or healthcare organizations Con actiently, the system is responsible for data loss.

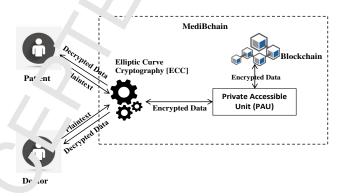


Figure 2: An application of MediBchain

Figure 2¹ depicts the design of our platform in which aforem ntioned problems have been addressed by storing the encrypted healthcare at in the cloud system. As a result, if our system somehow loses the control were clockchain, patients will be accountable for their data as they will control the encryption keys solely. Data sharing in our system is also being controlled by the patients. Vulnerabilities related to data preservation have been addressed in our system by using cryptographic functions along with blockchain technology. However, our system will store the encrypted personal data ensuring overal¹ privacy of the data such that even if system gets attacked by the attacker the stolendata vill make no sense to them. To get the plaintext of those encrypted personal data, attackers will require the keys. There is no identifier for these datasets, only encryption keys will be used to identify such encrypted and pseudonymouc² data.

1.1. Our Contribution

Our platform ensures that the private real moure data in cloud is controlled by only patient herself. The main idea of this work is to keep the sensitive healthcare data on the blockchain to attain account but v, integrity and security. Patients will have the overall control over the blockchain vhich their data will be stored. Present healthcare systems lack in pseudonymity as those only store the data in cloud, but our platform ensures the pseudonymity of patients. We achieve pseudonymity by using cryptographic function. Med 3chain will regain the interest of patients on EHR systems and will retain account ability, integrity, pseudonymity, security and privacy which are being lest verthe the increasing computational power of emerging technologies in EHR systems. Analysis of these attributes is discussed in section 3. Our contributions a east follows:

- 1. **Security and Arivacy guarantee:** The proposed platform guarantees accountability pseudonymity, authenticity and integrity along with data privacy.
- 2. **Analys.** ' digo ous analysis on security, privacy, accountability, pseudonymity and ' digrity nows how our platform achieves the above mentioned properties.

¹P₁ vate Acc 3ssible Unit (PAU) is the intermediary unit between blockchain and data sender or receive.

²Dseudonymity refers to the fact of using disguised identity.

An tysis of security terminologies are given in section 5.

3. **Evaluation:** We have implemented smart contract and show. If event analogies of costs (e.g., transaction cost, execution cost). Then we have evaluated a java implementation of input and output generation algorithm using Elliptic Curve Cryptography (ECC) for our system. Experimental results will help to compare several aspects of EHR system and will help to decide whether accept our platform or not.

Organization of the paper: The remainder of the paper is organized as follows: Section 2 describes the related work. In section 3 v e d scuss the preliminaries. In section 4, we describe our platform. In Section 5, we evaluate the platform and analyze it formally. We give some concluding remarks in Section 6.

2. Related Work

Some national level frameworks based on casted for electronic medical system have been proposed in [14, 13, 25]. Para et [25] proposed a model which is cloud-based and deals with patients' private data. This model ensures cost effectiveness, and this system was designed for rural areas where cost plays an immense role. Medical professior la and policy makers could serve the patients remotely through a cloud-based mood which stores all the imperative data in a single cloud. Patients were encouraged to share their data in the cloud so that they could get the medical service From the professionals remotely. Disease diagnosis and control could be in the ty this remote treatment. Data collection and data delivery are the key r sint in symptom analysis. Rolim et al. [27] proposed a framework where the system provesses data in the steps of data collection and data delivery. In this mode sensor play the role of collector which collects the data and sends directly to the system to store and work with this data further. These data would be accrossed by the medical professionals and sensors were proposed to be attached with the medical equipment in this system. Yin et al. [37] introduced cloud based parent centric system. This model includes three layers: data collection lay x, data management layer and data service layer. [21] described a blockchain base. ac ess control manager for heath data to enhance the interoperability of nis sy tem. Off blockchain mechanism with the involvement of public blockchai, was p oposed as an access control manager of healthcare data.

Controlla vility and Traceability are two key topics of privacy preserving systems. Also al. [35] proposed a model which is based on blockchain to help patients of the viscontrol and share their personal data easily and securely with privacy

preservation. This application based model also deals with Sec. Aulti-party Computing (MPC) and Indicator-Centric Schema (ICS). Simic et al. [23] showed a case study where the study concludes with the illustration of an nificant benefits of IoT and blockchain in a combined manner. In their work IoT cevices were proposed to be used as collectors of private health data of the patients', and real time data of patient could be saved in blockchain. Scal bility of the blockchain in case of Big data has also been tested in their study. Ek'law e al. [7] proposed a prototype named 'MedRec' which uses blockchair as a backbone and tried to find the security solutions for EHR systems. They rie a to give their prototypeintegrity, authenticity, auditability and data sharing 'brouz' olockchain. Elements of their system are: Registrar Contract (RC), Patient-Provider Relationship Contract (PPR), Summary Contract (SC), where RC n. ns t'e identification strings of the participants to their Ethereum addresses, PKN issues contracts between two nodes in the system when one node stores and manages medical records for the other, SC locates the participants medical record history. Jun et al. [3] proposed a web-based architecture where they showed a secured accessing multiple patient repository system. They concent red is ainly on lifetime repository of health data, which consists of client application (CA), central access-control (CAC), local access-control (LAC) and Hos real information system. Linn et al. [21] described a blockchain based access conu manager for health data to enhance the interoperability of this system

The backbone of our work is b'ockchain. Blockchain technology is popular for its application in Bite in c syptocurrency [26], which is a public ledger to hold and maintain the transpers all ata and integrity [31]. One of the reasons for using blockchain technology in cryptocurrency is its decentralized digital ledger property, which was presented by Nakamoto [22] in his Bitcoin cryptocurrency framework. Blockchain in data structure has been modeled by blocks which is linearly sequenced. Fach block contains the cryptographic hashes corresponding to the previour and current block to ensure continuity and immutability of the chain. Chaining riech anism ensures integrity of this secured data structure.

2.1. Bloc chain.

Figure? exhibits the structure of blocks in the blockchain network. In the figure each block is connected to its previous block by the hash of previous block. Blocks store he time-stamp of being mined in the network. Mining takes place in the network by solving mathematically complex problems. Miners compete each

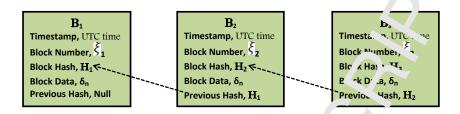


Figure 3: Structure of Blocks in bloc'schain

other to mine the block so that they could earn some eryptocurrency. In our platform miners will get Ether from Ethereum Network for unining, and our platforms Ethereum account will be charged against it. Simple Ether transfer functionality will be used to transfer the Ether from our account Each block contains corresponding block number and data that has been given to store in the blockchain which has been denoted as δ_n .

Blockchain-secured transaction-based technology [1] gives the users a better security. Bitcoin as well as blockchain as no been failed since these were introduced [6]. The network is shared and information is stored throughout the whole network, thus increasing the reliau." by on this technology. All the information is treated in a redundant way in blockchain [28]. Blockchain is distributed but it remains all the same for it' noc's ensuring the integrity [4, 34]. Centralized database can be corrupted an needs: third party to maintain it. To change the history of the blockchain any individual has to control at least 51% of the chain and it will cost a lot to challeng the impautability of blockchain. This immutable architecture [2, 30, 32] is a Flessin, it archival science too. Identities in the blockchain are covered by pseud n. ms by which privacy for the participants is ensured with a very high degree [15]. Cyptographic authentication of the time blocks with time-stamp allow the entire network to hold the logs for any interaction in the blockchain. Blocks ain ensures the verifiability of the users. Other than above discussed characte istics some author explicitly mention the key points like trust enabling notice [1, 2, 11, 33], Consensus, Transparency, Smart contract etc. Blockchair gives a distribution oriented service to be used as a storage. All the records that may be stored in the blockchain have to use smart contracts[16, 9]. Smart cont. ats determine the record of data and conditions in the blockchain. These contracts, as a form of code, give a huge power to the programmers to read and write ov r the blockchain [9]. As storage, blockchain provides accuracy and

reliability to it's users and protects the data from fraud and being tampered or

corrupted [18]. Blockchain as storage maintains proper decentral. at in and true redundancy, total privacy and cost reduction [10]. Decentralized web will be the future of this era[36].

3. Preliminaries

In this Section, we explain each properties (e.g., security, rivacy and management) that our protocol achieves. Finally, we introduce the building blocks of our protocol.

3.1. Properties

3.1.1. Security and Privacy

We briefly describe each of the security and privacy properties in the context of our system below.

- 1. Pseudonymity: No entity will be one identify any party of our system because users are being identified by a dynamic key. As a result users are keeping their selves pseudonymou. Data will not be identified by just seeing it.
- 2. Privacy: Only registered parties will be able to interact with the system. Even a registered party win not be able to access the private raw data of other parties.
- 3. Integrity: Authenticated parties will be able to store private data.
- 4. Accountability: Each back will be identified by corresponding block-id. Only authenticated parties will get them and will interact with them.
- 5. Security: P'rtie' will keep their encrypted data in the system which ensures secured environment for them.

3.1.2. Management

• User need a register once and by providing the ID and PWD ⁵ they can easly get it to the platform.

⁴Pe audonyn ty and anonymity are two different things. Anonymity refers to the fact of being unknow. in or *t* system users are identified with dynamic keys, hence users are pseudonymous.

- PAU will act as a Trusted Third Party (**TTP**) of our system, a it will be the medium between user and blockchain.
- In the case of Block id sharing, users need to be very care ?..' because untrusted access will make the platform vulnerable for u. particular user's data.

3.2. Cryptographic tools

Here, we describe Elliptic Curve Cryptography (ECC) [17] which has been used as the cryptographic tool to provide proper cryptographic functionality to the users. Formal definition of ECC will be given here.

Definition 1 (Elliptic Curve Cryptography) Enteric Curve Cryptographic scheme use the trapdoor function which means if we compute B from A through trapdoor function then it is mathematically infeasible to research A from B.

$$A \xrightarrow{\mathsf{trap}} \to B$$

All the functional properties of ECC are described:

Encryption Scheme:

Choose, Elliptic group $\mathbb{E}_p(a,b)$ and generator point, $\mathbb{G} \in \mathbb{E}_p(a,b)$ such that the smallest value of n for that $n\mathbb{G}$ \mathbb{G} a very large prime number.

Message, \mathcal{M} is encoded in to point $P_n \in \mathbb{E}_p(a,b)$

Both sender and receiver so lects $\sim r$ ivate key, $n_A < n$

compute public key P_A , such that $P_A = n_A \mathbb{G}$

Cipherte. voint,
$$P_C = [(\mathcal{KG}), (P_M + \mathcal{K}P_B)]$$

(K is the random integrand P_B is the public key of receiver here).

Decryption Scheme:

Plair expoint, $P_M \leftarrow (P_M + Kn_B \mathbb{G}) \leftarrow P_M + KP_B$ only receiver knowing private key n_B will retrieve this point, P_M by removing $n_B K \mathbb{G}$.

4. MediB Lain Pastocol

In this rection we present the architectural as well as the design view of our platform. **Table 1.** describes the notations that are used in the next sections.

Table 1: Terminology table

Notation	Description
ID	ID of the User
PWD	Password of the user
U_D	Encrypted user data
U_{id}	Block id, where user data will be aved
ID_X	ID of the User X
PWD_X	Password of the user X
U_{DX}	User X's Encrypted data
$\mathrm{U}_{\mathrm{idX}}$	Block number, where user X's a. 'a is saved
Secured channel	Obtained by the authentications process of our system
$\mathcal{T}(\delta_{\mathrm{n}})$	Transaction of δ_n through wart contract
$\mathcal{H}_{\mathcal{M}}$	Set of all identical haves
Γ	Address of the issuer
ν	Address of the me sage sender
$\delta_{ m n}$	Number of categories in the smart contract
$\{S,\!\mathcal{R}\}_{\mathtt{authenticated}}$	authenticated sender, S and receiver, R
$\{S,\!\mathcal{R}\}_{\mathtt{authenticated}}$	Unauthentic \mathcal{L}^{1} Dia tied, \mathcal{S} and \mathcal{R}
$\mathbb{B}_i, \xi_i \& \mathbb{H}_i$	Property of diffe. ant blocks

4.1. Overview of Our Protoco.

Fig. 4. shows the hig! level view of our platform. The following entities and their roles are described by each of the day here.

Data sender is the patient, who will send her personal healthcare data to the system. Data sender has since the vital role in case of data preservation. Data that will be sent to the system must be accurate otherwise wrong data will be detrimental for patient because the whole treatment depends on this sensitive data. However, our system will take the encrypted data from the users. Encryption of data will be done in the very comming of MediBchain's process execution.

Data rece. ver will request for the data after authenticating itself to the system.

Regis tration Unit will act as an authenticator. When any party (Sender or Receiver) will come for the first time to take the service of the system; it will store the ray and PWD to be used further. Each party will have to register for once

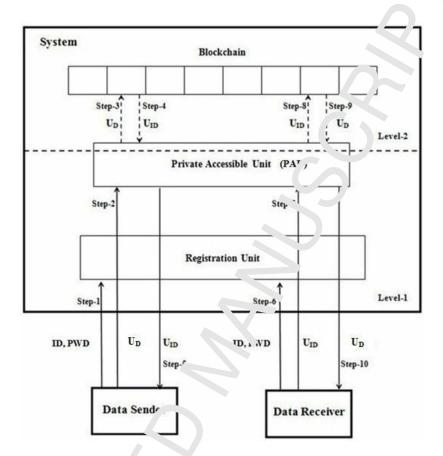


Figure i: High level view of this system.

and need to preserve and TD and PWD. Further they just have to log in and access through secured channel for transaction of their private data in the cloud.

Private Accessible Init (PAU) Both the parties of the system will be able to interact with FAU after authentication. It needs a secured channel to interact with PAU because it rugh this unit they will send their private data to the System. It is the intermy unit for both the levels of our system, through which the element of one level will a iteract with the other.

block chain vill hold the data of the users. Each transaction in the blockchain will return a lidentifier. Transaction identifiers will help the users to access the data further.

For better understanding our system is divided into two levels. Let al. is Graphical User Interface (GUI). User will interact with our system the rugh this level. Elements of level-1 are: Registration Unit and PAU. PAU is the element of both Levels so it will work between level 1 and 2. Level-2 is the broker do of our system, which interacts with low level elements of this system through PAU. Element of level-2 is: blockchain. blockchain is being used as a repository of healthcare data in our system. Our platform uses permissioned blockchain which will require authentication to access.

Steps in the system: Steps of our system could be defined from Fig. 4.

Step-1: Data sender will request with the ID and PWL to have access in the system.

Step-2: Upon accessing the system in step-2, Dan sender will send data to PAU for storing.

Step-3 & 4: Step 3 & 4 will take place in level-2 of our system, where PAU will send U_{ID} to blockchain and it will return U_{II} 101 future access to the blockchain and also for finding the exact Block v here to 9 data were saved.

Step-5: In this step PAU will return the Control to Data sender which was given by blockchain.

Step-6: From this step rest of the step are related to Data receiver. As step-1, this step also requires sign in r are sign in Data receiver can request for the data.

Step-7: In this step Data receiver vill request for the data to Private Accessible Unit along with the U_{ID} . AU will receive the U_{ID} for further use.

Step-8 & 9: Step 8 & 9 a. sar e as step 3 & 4 but the data are not same for this steps. In step-8 PAU vill request the blockchain along with the $U_{\rm ID}$ and in Step-9 blockchain will return it.

Step-10: This is ne final step where PAU send the private data to the Data receiver.

4.2. Formal Vesc iption of Protocol

In this section we will define how Data sender, Data receiver, and our system will work altoge her in case of sending and receiving the data. In case of data transmission in our system parties need to go through a step called registration. After confirmation of the Registration Unit that party can access the PAU.

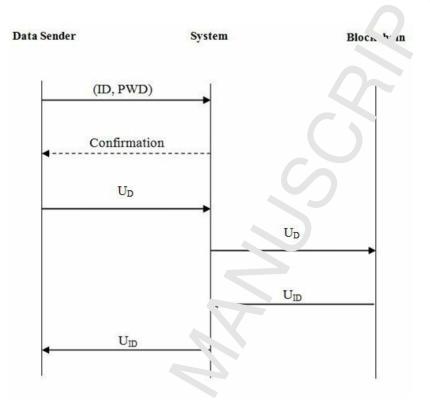


Figure 5: Low level view of Sending Protocol

4.2.1. Protocol between 'Jata sender and System:

Fig. 5. Shows the low x releview of sending protocol. A patient will play the role of a data sender in this protocol. Encrypted data will be sent to the system. Generation of ciphertexts collected depend upon a function known as encryption function. General zero form of this function is Enc(x,y). Below we will see how this function work.

$$Enc(key, Data) = U_D$$
 (1)

By providing κ_{ι} and the health data to this function data sender will get U_D and will send $\dot{\iota}_{\iota}$ to the system. Public key encryption technique (e.g., Elliptic Curve Cryptography (EUC)) will be applied to encrypt the private data.

Suppose X is D at a sender of our system. At first X will request for getting into the system D providing the D and D our system will send confirmation to D if D or D describes the right D and D or D or D to D through a secured

channel. Secured channel will provide the security to the transacies on of data . In this stage PAU will interact with blockchain and this interaction, with the blockchain will be done by the smart contracts of our system.

In our system smart contracts have been designed in a way such that blockchain will return the number of that block which has been denoted at U_{id} . Each block has a unique id which will work as an id of a specific patient. PAU will get the U_{id} on each transaction of data in the system for X it will be $U_{i'}$. PAU will send the U_{DX} to the blockchain then smart contract will return the special id U_{idX} , for X. After that PAU will send the U_{idX} to X and end the protocol. X has to store this U_{idX} otherwise next time X will not be able to access her personal private data. Getting the U_{idX} is the confirmation for X that means the data has been kept to the system and then X could log out and end the secured X nannel transmission with the system.

4.2.2. Protocol between Data Receiver and Sys. m:

Receiving in our system will take two larens of authorization. Because after registering or signing into our system parties will have to provide the U_{id} to get their data back through the secured change. In this phase if they fail to submit the U_{id} then they will not be able to write their data. U_{id} is the key to receive the actual data. Fig. 6. shows a low level view of receiving protocol.

Suppose user X wants to retrieve her data which she sent to the system in sending phase. As like sending phase this hase is also controlled with the authentication or Registration unit where Y has to sign in first then will be able to access our system. This sign in lequires the ID and PWD of the user which was given in the registration phase. If X raovides appropriate ID and PWD only then the system will send confirmation. After getting the confirmation X will be able to interact with the system unrough a secured channel. In this interaction with the system, X has to provide her U_{idX}. After getting the U_{idX} system will interact with blockchain. This interaction will take place in level-2 of our system. Only PAU can interact with blockchain, here the smart contracts of our system will be the medium.

Smart cor tract v 'll send the U_{idX} to blockchain for retrieving the data of X from it. 256 bit 'hash c' the corresponding block number will be checked in the smart contract, when the hash will be matched with any block then it will continue the proce 's to relieve the data. Otherwise this exception will be handled through the smart contracts.

Su, 'bo' - the hash of any block is,

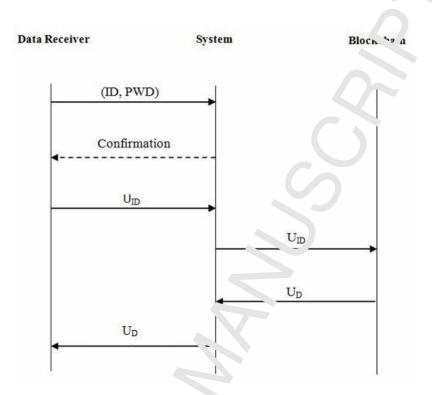


Figure 6: Low level view of Receiving Protocol

0xe3b1c14298fc1c149.fbf4c8196fb92427ae41e4649b934ca495991b785

Only if the hash of $U_{'X}$'s corresponding block is same then X will be able to get her data. In our system, blockchain will return the U_{DX} to PAU and it will be redirected to X lat x. After this data retrieval session will have it's end.

X will get her U_{DA} w' ich has to be decrypted to get the actual raw data to decrypt the data user n' ed to u > Dec(x,y) function.

$$Dec(key, U_D) = plaintext$$
 (2)

X will use equate n 2 with key and U_{DX} to retrieve the raw data.

4.2.3. storage of our system:

Our system will store the ID and PWD for authentication and response purpose. Our system solely will manage these private data in the cloud without depending on any other trusted third party (TTP) apart from the PAU. Each time

when user will store the data she will get a new block to write so t_i . b'ock-id will change by time. ID and PWD is dependent on party but U_{id} is dy. and with each data storing process.

4.3. Programmatic view of MediBchain

Smart contract of our system has been presented in this pape, through some algorithms. These algorithms have been designed to be concerted in any blockchain based language (e.g., Solidity, Golang). Contracts of our system are written in Solidity language and all the results of this paper are also be ed on Solidity based environment. Algorithms 1-3 will be appropriate for any environment designed for blockchain environment.

Algorithm-1 describes how our system will check the insuers verifiability. All the hash of our system is denoted by $\mathcal{H}_{\mathcal{M}}$ and all the tailed \mathcal{H}_{i} must be a part of $\mathcal{H}_{\mathcal{M}}$. Here, i refers to particular number of \mathcal{H}_{i} .

$$\mathcal{H}_{\mathcal{M}} \leftarrow \{\mathcal{H}_1, \mathcal{H}_1, \mathcal{F}_{13}, \dots, \mathcal{H}_i\}$$

 $\mathcal{H}_{\mathcal{M}}$ is the set of all identical hashes of our system that will be provided in the time of account creation in the blooking network. $\Gamma \& \nu$ are part of $\mathcal{H}_{\mathcal{M}}$ and play significant role in transaction. Two different notations have been used to reduce the complexity of Algorithm 1, issuer of contract has been denoted with Γ and data uploader/downloads has been denoted with ν . Here, issuer is the address who runs the contract and message lender is the address who sends the message. If both of them are not so not fine Algorithm-1 will return false.

Algorithm 1: Checking on 1ss aer and Sender

The Ligorithm is important for security and accountability of data transaction.

It will work in between the time of smart contract execution and the data transaction (e.g., upload, download) between MediBchain and blocker. In. L. vesdroppers could take a chance of data manipulation in the meantime. All use accounts of this system will be the part of \mathcal{H}_M and also the initiator of contract and data uploader/downloader will be same. Execution of rest of the contract will be dependent on the similarity of $\Gamma \& \nu$. Algorithm-2 will be in tiated after algorithm-1, in which δ_n represents the number of categories to be held by the structure of data in our contracts.

Algorithm 2: Transaction of Data

```
Result: Data Upload
 1 struct Data \leftarrow \sum_{1}^{n} \delta_{n}
 2 Data[] data;
 3 bool←0;
 4 while n do
 5

    getting input from message sender,v

        if \nu returns string then
 6
             data \longleftarrow \sum_{1}^{n} \delta_{n};
 7
             bool \leftarrow 1;
 8
             return bool;
 9
        else
10
             return bool;
11
        end
12
13 end
```

Algorithm-2 will be created after fulfilling the conditions below. Iff,

$$\Gamma$$
 && $\nu \in \mathcal{H}_M$ and,
 Γ (1ssuer) = ν (message sender)

Here, Γ is the raddres, who runs the contract and ν is the address who sends the message. If by the rather are not same then this algorithm will return false. Users' (patients') data vill be having different categories to be inputted. Different categories mean that the healthcare data come in different types, suppose user wants to save B. bod subars data and also Blood pressures data these two are different. By category we refer to this scenario that the user can store different diagnostic results in a back. Hence, we have designed two different contracts. In Algorithm-2 each structure will hold maximum four different types of healthcare data to be steed the block if we change data part as follows,

$$data \longleftarrow \sum_{n=1}^{4} \delta_n$$

We have another smart contract which takes maximum eight a rerent types of health data to be stored in the block. For that we need to charge data part again. So above data part will be changed as follows-

$$data \longleftarrow \sum_{n=1}^{8} \delta_n$$

We have shown some computational analysis in subsect on- 3.3 using the variation of data storing capabilities of different smart contrac.3.

Line-1 is showing that the structure of smart contract can take n number of individual data from a particular patient at a time. In the loc p da a will be assigned to its corresponding structure in line-7 and then the a to will be written in the block in the same contract. A particular structure will be written in a particular block. As mentioned earlier each block of blockchain hold different id which is not same as \mathcal{H}_M . \mathcal{H}_M represents the account id of blockchain network whereas hash ids has been denoted with H_i . A *bool* variable was been returned from Algorithm-2 as a flag for Algorithm-3. In Algorithm-3. \mathcal{E}_i represents the block number and H_i represents the hash of particular block.

Algorithm 3: Block-id Generation

```
Result: Block-id
 1 \xi_i, \mathbb{H}_i;
                         ▶ Will hold in \mathbb{H}_i \leftarrow Hash of Block, \xi_i \leftarrow block number
 2
 3
   while bool do
                                                           ▶ Returned value from Algorithm 2
 4
        if bool \leftarrow 1 ther
 5
              \xi_i \leftarrow block Nu. ber);
 6
              \mathbb{H}_i = \text{bloc! } . \text{`-lockhash}(\xi_i);
 7
              return Hi
 8
         else
 9
              return 1 '''1;
10
        end
11
12 end
```

Algorium-3 will return hash id \mathbb{H}_i if all the requirements will be fulfilled by the co. tract. I will take a variable named *bool* by which this algorithm will define them. To return block-id, \mathbb{H}_i or not. Functions block.Number() and block.blockhash() are the syntex of Solidity language, where block.Number() will return the corresponding block number ξ_i and block.blockhash() will return \mathbb{H}_i

$$\mathbb{B}_i \ni \xi_i, \mathbb{H}_i$$

 ξ_i and \mathbb{H}_i are the properties of each block, \mathbb{B}_i by which our system will work

$$\mathbb{H}_i \leftarrow block.blockHash(\xi_i)$$

Programmatically each \mathbb{H}_i will be generated from it's corresponding ξ_i . As instance, if block.blockhash() gets ξ_1 as a parameter it vill return \mathbb{H}_1 or if it gets ξ_{20} the function will return \mathbb{H}_{20} . So the relation can be written as,

$$\{\mathbb{H}_1, \mathbb{H}_2, \mathbb{H}_3, \mathbb{H}_4, \dots, \mathbb{H}_i\} \equiv \{\xi_1, \xi_2, \xi_3, \xi_1, \dots, \xi_i\}$$

5. Protocol Analysis & Evaluation

- 5.1. Security Analysis
 - **Pseudonymity**: Data Sender, S and Fecciver, R will not be identified by any party during transaction.
 - Pseudonymity of S: After any entication S will upload the encrypted private data, U_D . Any outer party will not be able to identify S by looking her U_D because of it's identificationless attribute.
 - Pseudonymity of \mathcal{R} : \mathbb{H}_i will be used to trace particular \mathbb{B}_i of the blockchain which holds the private data of \mathcal{S} . During transaction \mathcal{T} party will hold the \mathbb{H}_i to have her U_D back from the system, these \mathbb{H}_i s are as sensitive at the private data for receiver. \mathbb{H}_i will be held by only our party which ensures the pseudonymity of Data Receiver because no one will be able to detect \mathcal{S} during \mathcal{T} or even after \mathcal{T} because of encrypted property of U_D .

Suppoor $x\{ID,PWD\}$ is the function for authentication,

$$\alpha\{\text{ID,PWD}\} \longrightarrow \{\mathcal{S},\mathcal{R}\}_{\text{authenticated}}$$

- **Privacy**: Registration Unit and U_D ensures the privacy of the $\{S,R\}_{authenticated}$ and data respectively.
 - P.ivacy from system: Parties, $\{S,R\}_{\text{authenticated}}$ of our system have privacy as pseudonymity of users is maintained. $\alpha\{\text{ID,PWD}\}$ will ensure

the access in the system. This controlled access of $\{\mathcal{L},\mathcal{R}\}$ uthenticated provide privacy to the users of our system. Therefore, \mathcal{L} of $\{\mathcal{L},\mathcal{R}\}$ authenticated can not be compromised any way.

• Privacy from other parties: S will have her dedica. $A^{\dagger} J_i$ in the blockchain to store U_D . So, if any $\{S, \mathcal{R}\}_{\text{authenticated}}$ of our system tries to access any other party's data it will not be able to access the particular block as each party will have their dedicated H_i .

Clearly, the former analysis guarantees a very strong rivacy of parties because only $\{S,R\}_{\text{authenticated}}$ will be able to cross as well as retrieve data from that particular \mathbb{B}_i .

• Integrity:

- Access request data integrity Each time S or R tries to access the system, she needs to authenticate burself primarily. This access request needs to be done by both the dynamic entities. S and R of system. These access requests will require correct ID and PWD, which will be generated by party itself and will be holding by the database of system. So without S or R and system these authentication data will not be known by anyone. Ly which system guarantees the access request data integrity.
- User data Integrity Use of Enc(x,y) function ensures the data integrity as the data in 'be blockchain will make no sense to any other person except the data owner. After retrieving the data from the system $\{S,R\}_{authentica}$ need to decrypt the U_D with Dec(x,y) function. In order to eak this integrity level attacker needs to break the security of unearly negencryption scheme, ECC.

All the cata that are related to our healthcare data management system guarantees in the crity.

o Accumtability:

• I ansactional \mathbb{B}_i : When any party will come to save it's data to the system a unique number or nonce, \mathbb{H}_i will be returned which leverages the accountability of our system. Only party itself will be holding this

nonce which makes the party accountable for it's $U_D \ \ \ U_{\text{\tiny C}}$ 'a' se without valid information about

 $\mathbb{B}_i \ni \xi_i, \mathbb{H}_i$ party will not be able to access her $_{\mathbf{F}}$ ivate data from blockchain.

- PAU as bridge: Interaction of $\{S,R\}_{authentic}$ and which the system is controlled. This controlled path refers to the secure 1 channel which will be created by the party itself through $\alpha\{I,PW,D\}$. Through this channel $\{S,R\}_{authenticated}$ will interact with $P^{A,T}$ which is a bridge between the system and blockchain. Secure it charmel makes the bridge accountable for secured T with blockchain.
- Security: Each \mathbb{B}_i will be dedicated to $\{\mathcal{CR}\}_{\text{authenticated}}$ and their \mathbb{H}_i is secured as integrity is guaranteed in our platform. As a result, these \mathbb{B}_i will not be accessed by any $\{S,\mathcal{R}\}_{\text{authenticated}}$. If attacker somehow manages to intrude into the blockchain network patients' sensitive data will make no sense because of encrypted "ribus" of data. Accessing the raw data of patient will need the keys and $De_{\mathcal{C}(A,\mathcal{C})}$ will return the raw data to parties. So, the data security is guarance. "in our platform.

The equation for **Transaction**

$$\mathcal{T}(\delta_{n}) \longleftarrow \left\{ \left\{ \forall \mathcal{H}_{\mathcal{M}} : \Gamma \in \mathcal{H}_{\mathcal{M}}, " \in \mathcal{H}_{\mathcal{M}}, \Gamma = \nu \right\} \text{ and } \left\{ \forall \alpha_{\{\mathcal{S},\mathcal{R}\}_{\text{authenticated}}} \{ID, PWD\} \right\} \right\}$$

After analyzing each of t'e p operties we can conclude with saying that no platform secures blockchain but d preudonymous healthcare data other than our platform-'MediBchain', in the 1 est of our knowledge.

5.2. Computation Tvaluation

We setup an Cryi onment to evaluate our protocol by writing programs using Solidity 0.4.11 and JATA 1.8 with a computer Intel(R) Core(TM) i5, CPU-3.30 GHz, 8 GB of RAM, Win 8, 64-bit OS. We deployed Elliptic Curve Cryptography (ECC) for generating, and retrieving the input and output respectively.

5.3. Paul sharing:

We test the computation time to generate the cipher texts. Each encryption is an isolated process. Fig. 7. presents the data encryption time versus string size of nealthcare data. We take several inputs to see how the rate of growth of

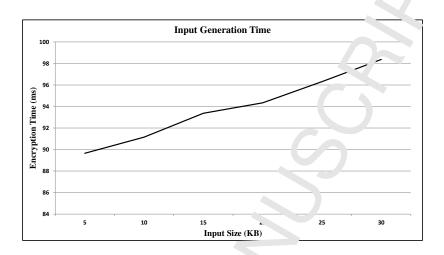


Figure 7: Computation in generating input

time for encryption changes with variable input size. We take 5 to 30 kilobytes of data to analyze the encryption time of different data size. From the resultant graph we can observe the rate of growth of curve is nearly linear which means the encryption time increases with increase of data size. Data sharing phase of our system is variable and independent process, variable means that input size could vary for different users and independent on each other.

5.3.1. Data manipulat on will mart contract:

The issues that have then mentioned in the manuscript could be solved with other technologies that through blockchain environment we get the proper distributive attribute which lacks in others. Blockchain gives us the option to use it as distributed leaght which makes the technology a viable option. Ethereum environment has been used to analyze the effectiveness of this new idea of EHR system over which we operating system. Ethereum is the most effective platform to run Darps (Pistributed Apps) using solidity language that is the reason why Ethereum platfor has been used to access blockchain.

Before stting access of a block in the blockchain network data will be accessed by our smart contract. Use of smart contract will cost some gas which is known as the cryptofical of Ethereum Virtual Machine (EVM). To run any dapp (distributed application) on the Ethereum environment the executed application will need to

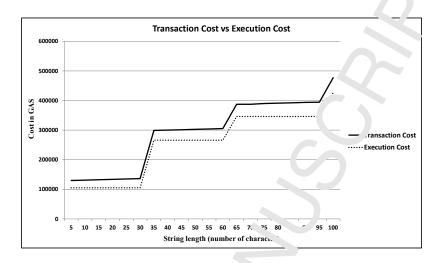


Figure 8: Computation cost of transaction and vecution of smart contract

have some transactions in the networ, in return of transaction the environment costs the executor some gas. Initiator of executor of transaction will get the gas in exchange of Ether in Ethereum exprominent. We evaluate two smart contracts, one with 4 inputs category other with 8 inputs category. In context of programming language which is number of variables to take input from party. Subsections 5.3.2 and 5.3.3 will depict the analogy of different terms of smart contract with 4 inputs category and 5.3.4 and 5.3.5 will depict the analogy between two different smart contracts with variable inputs, where . We tried to show some analogy based on the transaction and execution cost of our smart contract.

5.3.2. Transaction Cost vs Execution Cost:

Fig. 8. depict. the analogy between transaction and execution cost of smart contract. To have an accurate analyzing result we run the smart contract with different input sizes that varies from 5 to 100 characters of string. Curves in Fig. 8. shows the cost at ancreasing with the input size. But the rate of growth of these two curves is san a between the intervals and linear too.

5.3.3. Block-id generation costs:

O've of the key terms to be ensured while writing smart contracts was block-id generation. Block-id generation will cost for execution and transaction. We analy to block-id generation cost with different string length, but interestingly it

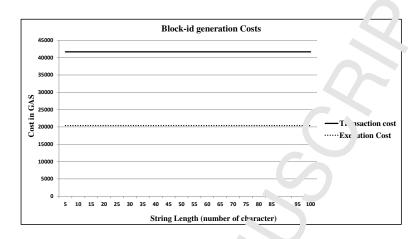


Figure 9: Computation cost of transaction and Cacution of smart contract

costs same for all the inputs. Fig. 9. shows the curves of execution and transaction cost of block-id generation. It is clear that each parameter is almost constant with the increase of the size of string. Transaction and execution cost is same for growing input size.

5.3.4. Transaction cost of variable inpu.::

Parties of our system may have to upload a vast amount of data in different categories. Smart contract may have to be redesigned so that we analyze the cost to see how our platform reacts with an increasing amount of category to store it in blockchain. Before this subsection we were talking only about smart contract having 4 categories to fake is input, but for having an effective analogy we will give 8 categories as input to see how the behavior changes of our platform. Fig. 10. shows us the analogy to tween two smart contracts in which one will take 4 inputs and other will ake 8 inputs. In Fig. 10, we can see that smart contract having 8 categories of input will cost higher, but the rate of growth of curves are similar and the cost with increase with string size.

5.3.5. Execution of tof variable inputs:

Fig. 11. presents the execution cost of smart contract with variable input. As explained those mart contracts may vary in different scenario, so that we present the execution costs' analogy in Fig. 11. The rate of growth of curves is similar but smart contract with 8 inputs will cost more gas while execution with increasing string lenguls.

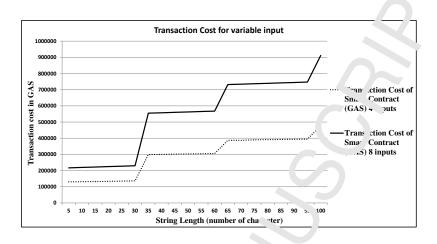


Figure 10: Transaction cost of smart contract with variable input

5.4. Output generation:

To get the plaintext or actual private her theore data of patient the data from blockchain need to be decrypted. As here are yption, decryption or output generation process is also isolated. All the output generation for the parties is independent from each other. To analyze the output retrieval time we take different sets of string 5 to 30 kilobytes of data at a single input to get an actual idea of output retrieval time for the patients. In 1 g. 12. curve shows that the rate of growth of time is related with the input size as the time is increasing for decryption with input size. The curve is rearly linear. Time is in millisecond in the graph that is computed with Java derive decryption. Elliptic Curve Cryptography (ECC) is used to generate the plaintext.

5.4.1. Input generation vs Output retrieval:

Generation of input and output is independent from each other. Encryption will take place in use time of giving input and decryption will take place in the time of output Fig. 13. depicts that two processes take very different amount of time while processing. With the string length both the time increase but encryption needs more than decryption. For encryption it takes 80 to 90 milliseconds where decryption needs less than 10 milliseconds.

6. Cenclusian

The puper presented privacy preserving platform for healthcare data in cloud. We have defined a set of security and privacy requirements for healthcare data

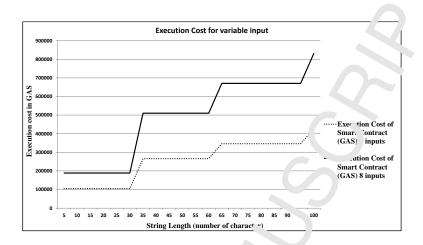


Figure 11: Execution cost of smar, ontract with variable input

management systems and argued why such attributes are necessary for a health-care data management system in closed. Our analysis shows that our platform satisfies all such requirements. Experimental performance evaluation shows that this platform runs well in blockchair environment. In the future we will try to explore the interoperability between different entities (e.g., diagnostic center, hospital, doctors, patients) of herence process, and another direction would be to address the issue of handling bey-the ft/loss mechanisms or key distribution techniques.

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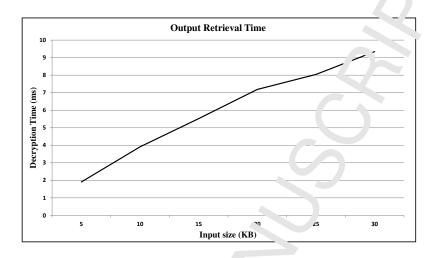


Figure 12: Computation time in generating actual output

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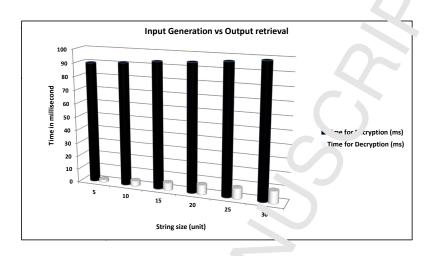


Figure 13: Input generation vs vur training of system

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Highlights:

- 1. We proposed a user centric EHR systems for Healthcare, which gives the total controlling power of the data to the users.
- 2. In generic EHR platform, it becomes easier target of intruders to not ide the system than totally breaking the security of the system. We solved this problem by implementing permissioned Blockchain along with the cryptographic function.
- 3. We explored the archival use of Blockchain in our platform by storing the data of users in the blocks of the permissioned Blockchain.
- 4. Controlling the pseudonymity of the users is a big challenge. Ver solved the pseudonymity issue by applying cryptographic function. We used Elliptic Curve Cryptography (ECC) to make the data safe from other party in this distributed ledge. System.